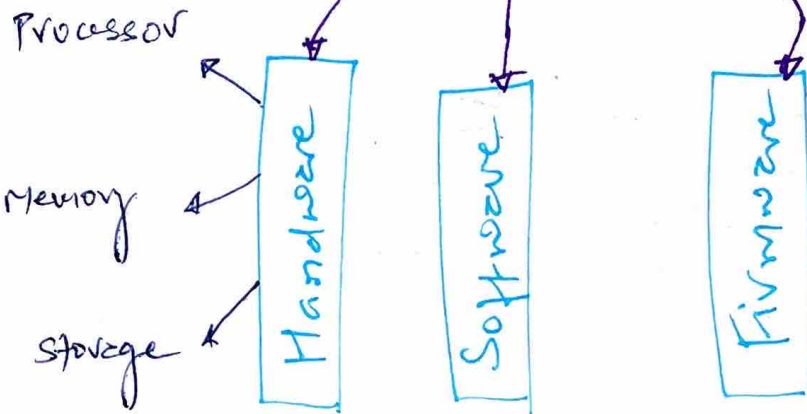


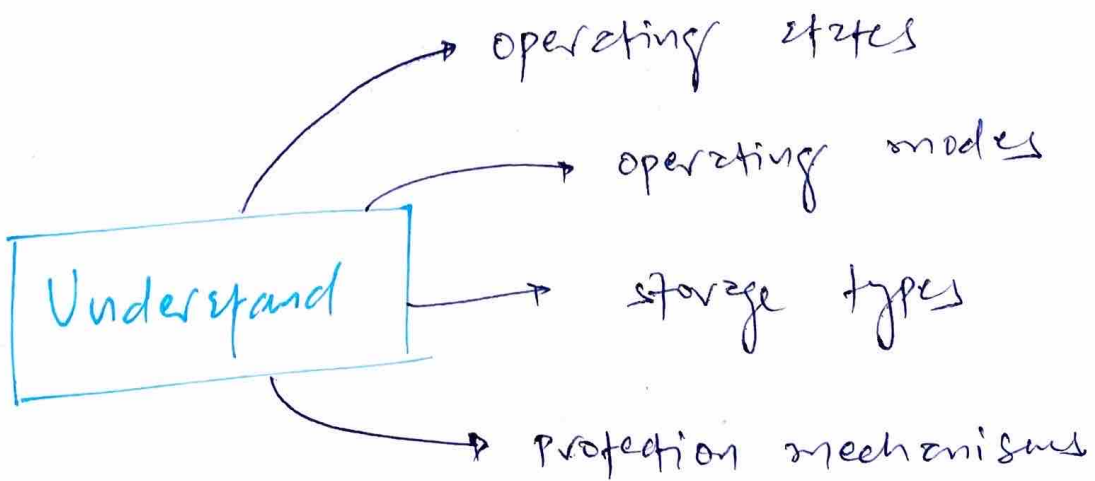
9. SECURITY VULNERABILITIES, THREATS, AND COUNTERMEASURES.

DEEP DIVE INTO COMPUTER ARCHITECTURE

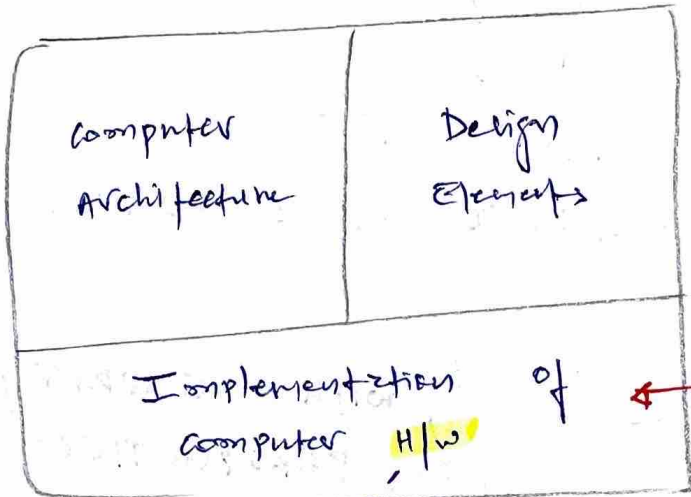
with SECURITY PERSPECTIVE



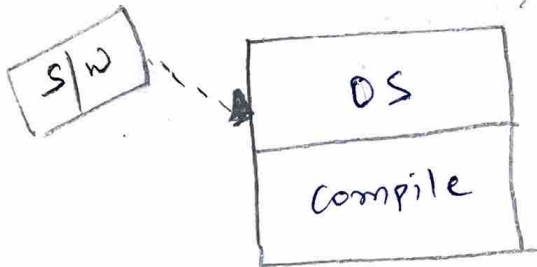
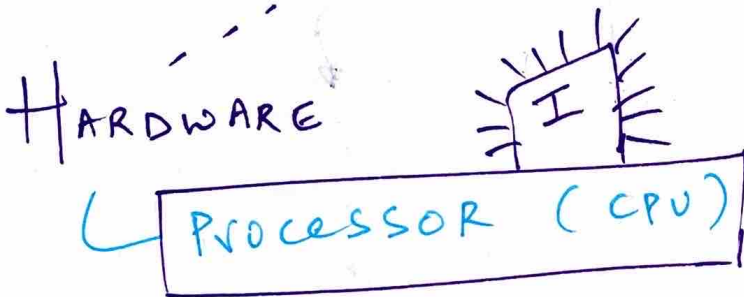
CHAPTER PERSPECTIVE



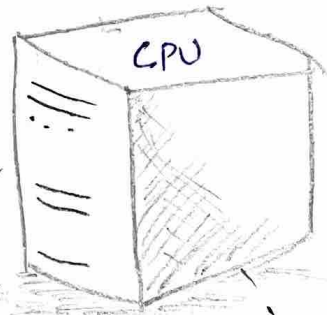
ASSESS & MITIGATE SECURITY VULNERABILITIES



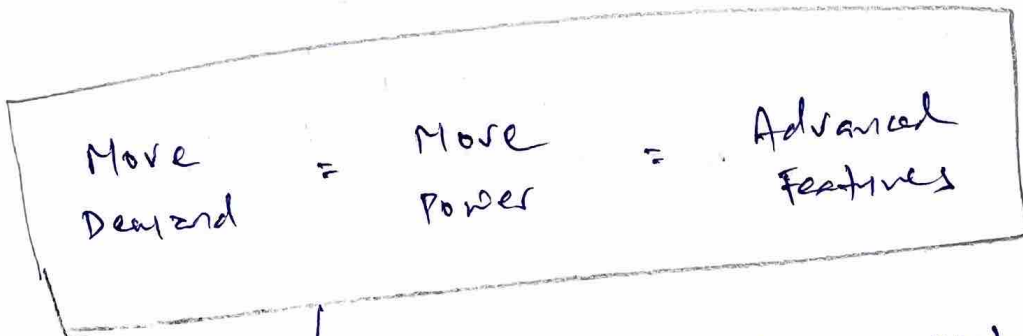
Foundational Understanding of CISSP



Translate s/w's
High-level programming language to Assembly language so, CPU can understand



Computational + Logical operations



Execution types (P.T.O)

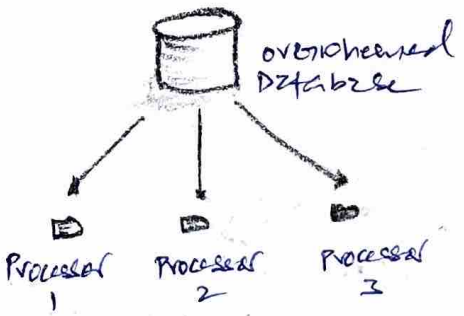
Execution Types

Multitasking

Simultaneous execution of one or more task / applications managed by OS.

Multiprocessing

More than one processor to increase computing power.



2 types

SMP

- Symmetric multiprocessing
- simple operation at extremely high rates
- Processors share common OS + data bus + memory

MPP

- Massively parallel processing
- For large, complex, computationally intensive tasks
- Thousands of processors where tasks are further break into mini tasks, distributed to other processors & decompose back again.

Uses multithreading at OS level

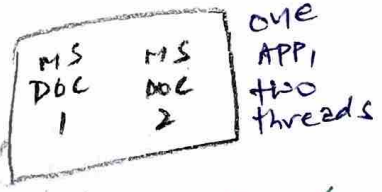
SMP send one thread to one process for simultaneous execution

when we open 10 word apps - each doc is thread, 10 threads but one process

multiprogramming

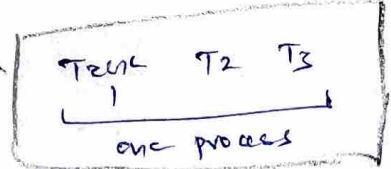
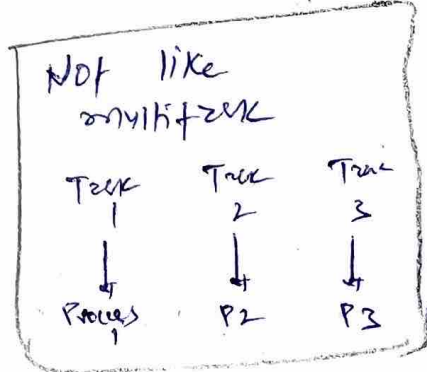
Similar to multitasking but takes place in mainframe system & require specific programming

multitasking is coordinated by OS while multiprogramming requires specific written software.



Multithreading

Allows multiple concurrent tasks to be performed within a single process



- context switching b/w active process reduce overhead, increase efficiency.

Processing Types



Design in such way
= Doesn't disclose information
to unauthorized clients

Other option is
viz policy. But
we have focus
on hardware
processor level

Address this at
PROCESSOR LEVEL

2 different
ways

Through policy
mechanism
to manage info.
at different levels

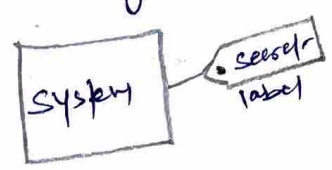
Through Hardware
(Higher security)

SINGLE STATE SYSTEM

MULTI STATE SYSTEM

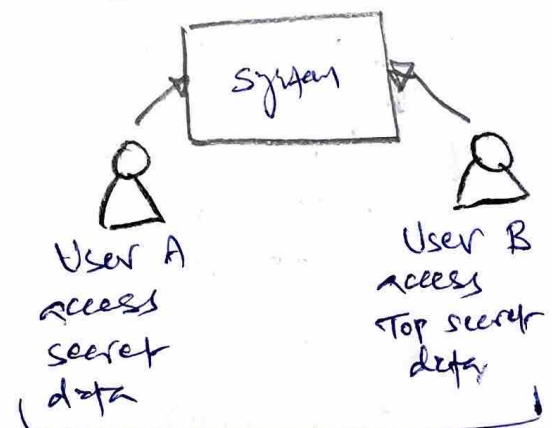
System admin approves
processor + system to handle
one security level at a time.

These systems are certified
to handle multiple
security levels simultaneously.



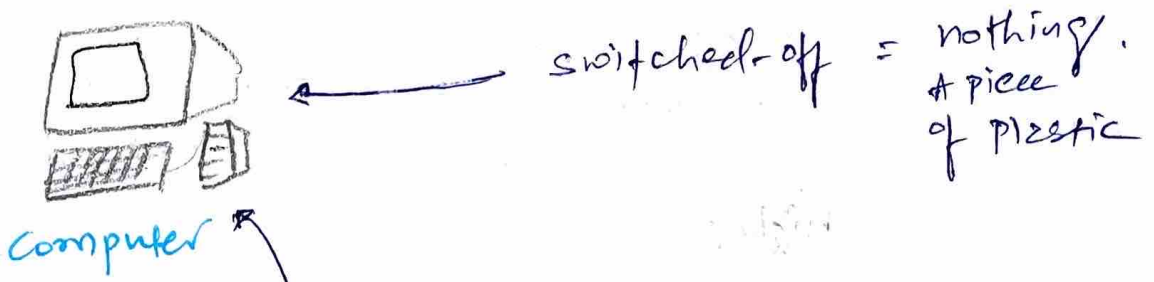
All users approved to handle
secret information. This
takes burden off to handle info.

- from Administrators
- Hardware
- OS



Technical mechanism prevents
information crossing b/w
two users (between two
security levels)

Protection Mechanisms

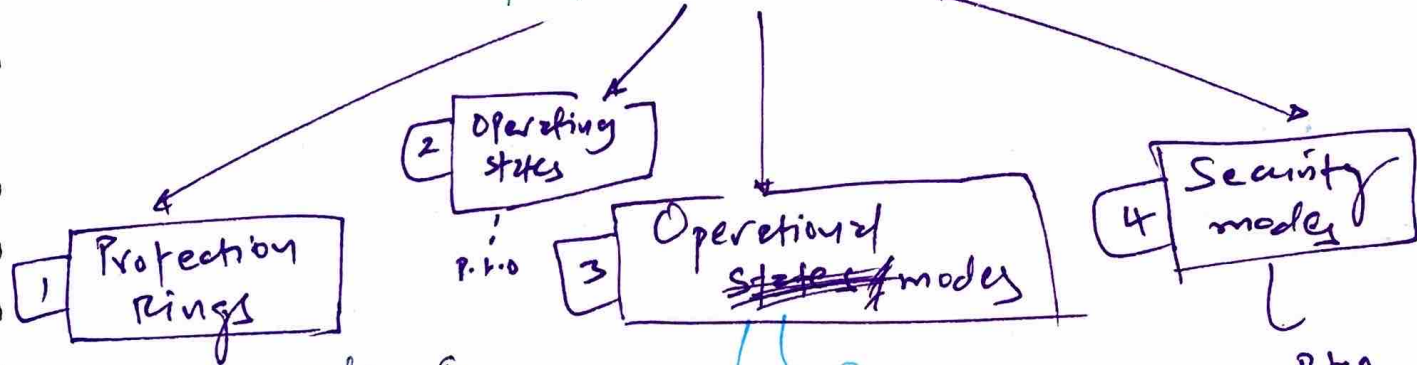


Computer

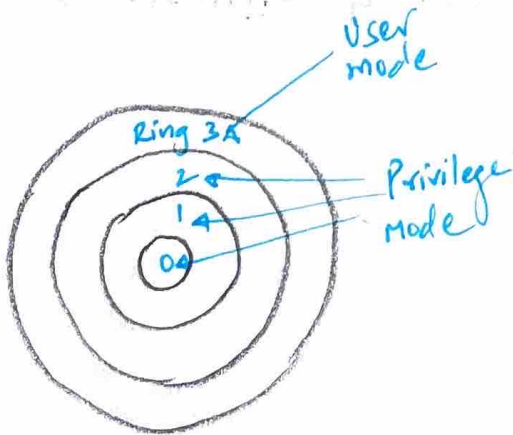
ON

Running computer has to address information security

Various protection mechanisms



1 Protection Rings
- organize code & components in OS into concentric rings.



0-2: Protection ring segregates OS into kernel, components & drivers

3: This ring runs applications & programs. P.T.O

① User mode
CPU
- Access to small set of instructions
- often executed in controlled environment such as VM.
- Prevents user's unintentional action + malicious user's intentions

② Privilege mode
- OS access to full range of instruction supported by CPU

- wide range of permissions, Be careful when you give privilege access.

P.T.O
X2

--- Protection rings (contd---) | 1

Ring 0 is **KERNEL**

Lower numbered Ring = **VIP**
 Ring 0 can access any resource / memory location. To move access & interaction with OS

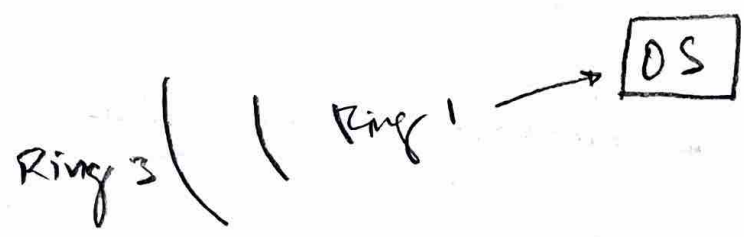
Ring 1 = OS components

Ring 2 = Drivers, protocols

Higher numbered Ring = ordinary → less access

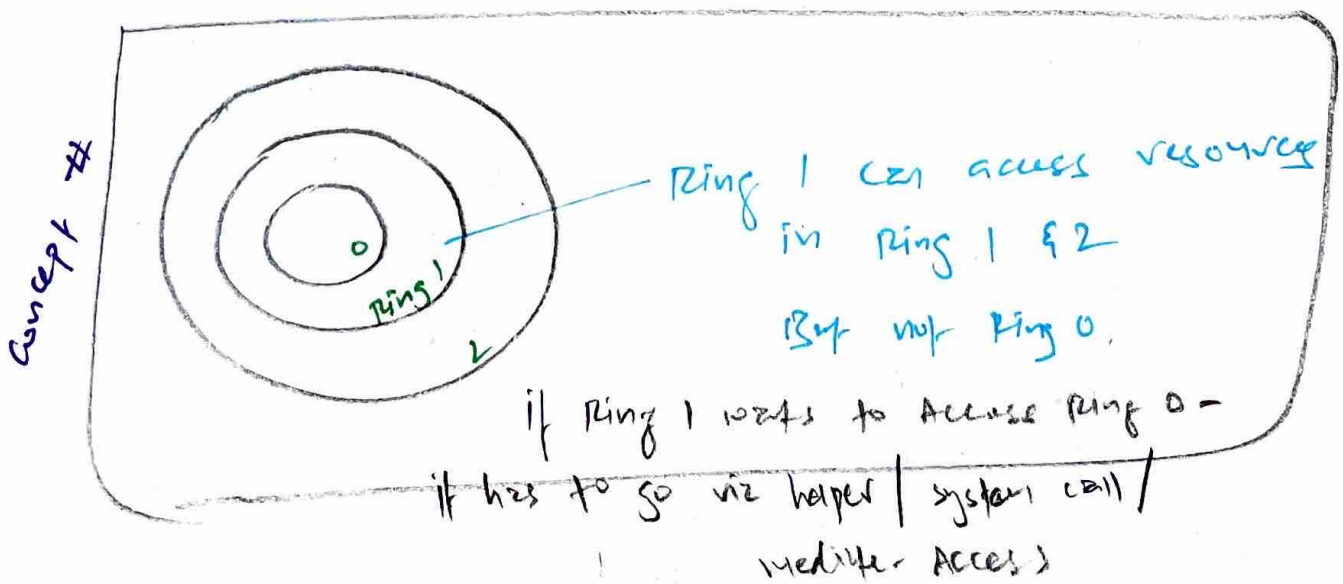
Mediated-access model (system call)

Ring 3 - User-level programs & Applications



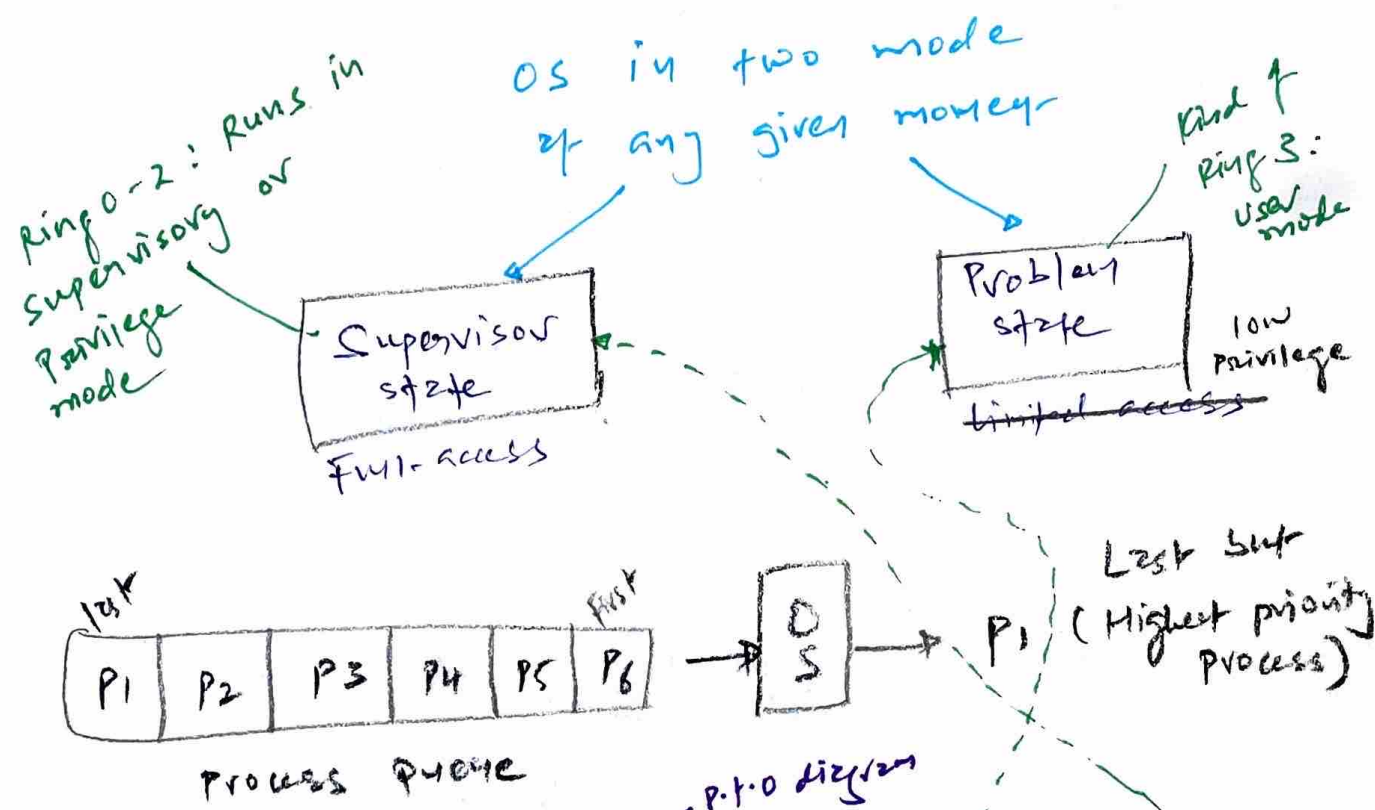
Ring 3 uses Ring 1 for helper / driver to access services in OS. Higher numbered rings need helper

Concept # Ring model → OS to protect & insulate itself from users & applications



2 Process states --- side concepts / part of protection using operating states

Various forms of execution in which process may run.



STATES

Ready (Queue)
Process is ready to run. it has memory + resource

Running
Process continue to run until time slice expires or process is blocked. In these cases, process goes to ready state (Queue) if time slice expires

Supervisory
when process wants to perform privilege actions

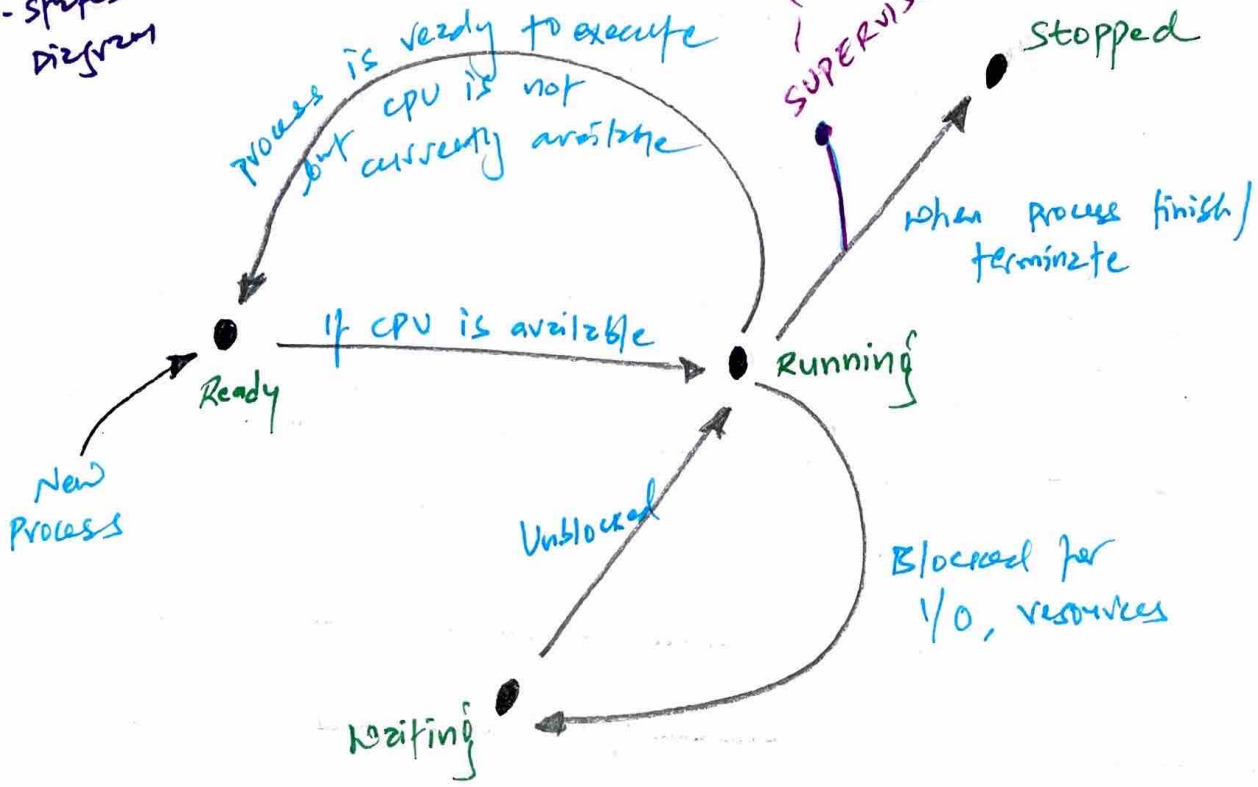
Waiting
Process waiting resource or device or access request

Process goes to waiting state (Queue) if block while waiting for resource become available.

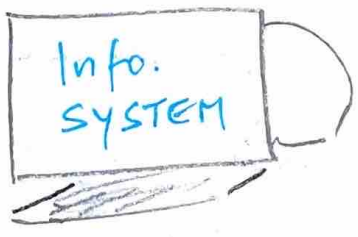
Stopped
Process is stopped when job is done or terminated. All resource + memory released so that OS can use for other processes.

4 Security Modes

Process states diagram



[Process scheduler Diagram]



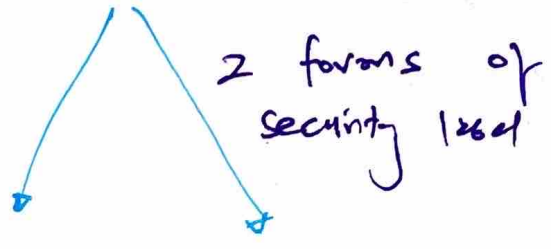
To Protect

Four Security Modes

- (Strong) Dedicated Mode**
 - Similar to single state system
 - valid need to know
 - security clearance
 - Access approval
 - (less strong) System High Mode**
 - Need to know not necessary for all information processed by system.
 - (weak) Compartmented Mode**
 - No need for access approval for all information in system (need for any)
 - P.T.O CNU
 - (weakest) Multilevel Mode**
 - Not everyone has security clearance.
- diff: 3 requirements

CMW: Compartmented mode workstations

special mode implementation



Sensitivity levels

Which objects must be protected

Information levels

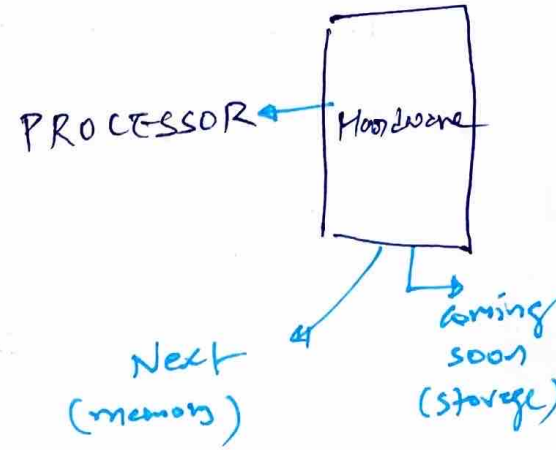
Prevents data overclassification

Table - Comparing Security model

Mode	Security clearance	Need to know	PDMCL (if CMW is used)
Dedicated	Same	None	None
System High	Same	Yes	None
Compartmented	Same	Yes	Yes
Multilevel	Different	Yes	Yes

II MEMORY

So far we learned



can't be modified / No writing allowed
ROM
 (Read-only memory)

Programmable ROM (PROM)

- can be modified to some extent

Erasable Programmable ROM (EPROM)

Ultraviolet EPROM (UEPROM)

- Erase with light
- End user can burn new information

EEPROM must be fully erased to be rewritten while Mask memory can be erased & written in blocks.

Electronic Erasable EPROM (EEPROM)

FLASH MEMORY!
 - nonvolatile storage for electronic erase & writing

Use of electric voltage to erase data

RAM
 (Random Access Memory)

volatile content = Temporary usage
 Power-off = gone

Real memory

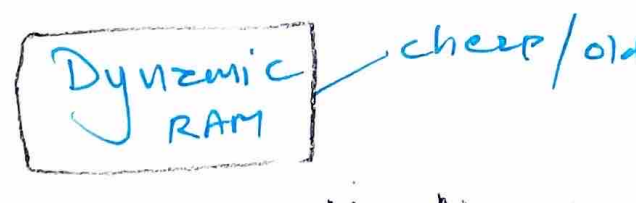
- Primary / main memory
- largest storage on PC

Cache memory

- takes data from slower device & store to faster device as cache to boost performance.



vs



- Uses flip-flop as ~~uses semi~~ logical device
- maintains contents unaltered as long as power is supplied

- Uses series of capacitors to store data

static / dynamic = Volatile

Control

Registers



- CPU's limited amount of onboarded memory
- Provide direct access to brain of CPU = ALU (Arithmetic logical unit)
- Advantage: this type of memory (register) is part of ALU itself

Memory Addressing

- For processors to refer various locations of the memory

- ↳ Register Addressing: Still memory directly into the CPU
- ↳ Immediate Addressing: Not memory addressing scheme but ~~more~~ how data is referred to CPU as part of instructions
- ↳ Direct Addressing: CPU is provided with actual address of memory location to access.
- ↳ Indirect Addressing: Memory address contains another memory address
- ↳ Base + offset Addressing

Secondary Memory



CD/DVD



USB



External drive

Special kind

- Storage device that contains data not immediately available to CPU
- Inexpensive compare to primary memory, holds massive amount of information

→ **Virtual memory**: Such as **page file**, OS uses for memory mgmt

→ Inexpensive but ~~slow~~ makes system paying operations slow

occurs when data is exchanged b/w primary & sec. memory are ~~slow~~ ^{slow}

Memory Security issues

FIRST UNDERSTAND

→ where data is stored? (what kind of memory)

→ How it is stored?

If memory device store = sensitive data

↳ **PURGE**

(before they leave organization)

cyber sketch

COLD BOOT ATTACK

Freeze memory chips

ATTACKS ON

memory image dumps or system crash dump to extract encryption keys.

Concern: Who will access data stored in memory while a computer in use

Data-in-use

Use this principle

OSG P 382

PROCESS ISOLATION

Process don't have read / write access to memory spaces that are not allocated to them.

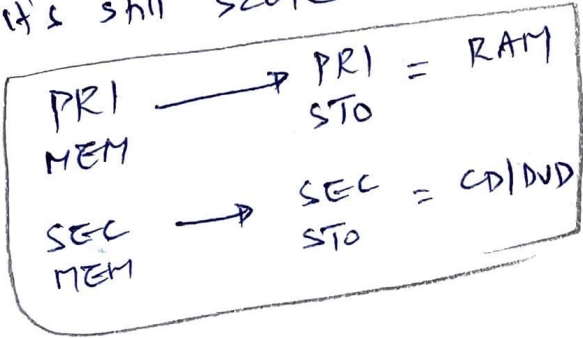
Implement VM per user / per process basis



Terms

Primary vs. Secondary

- Point compare with primary & secondary memory but it's still same



(Faster) — Random vs. Sequential (Slower)

~~Power-off~~ → ~~Data loss~~ = ~~Vol~~

Random = ^{+secondary} Primary storage device, allows OS to read/write anywhere / any point

Sequential = magnetic type = read/write data in sequential order

Hold massive storage = good for backup

Volatile vs. Non-Volatile

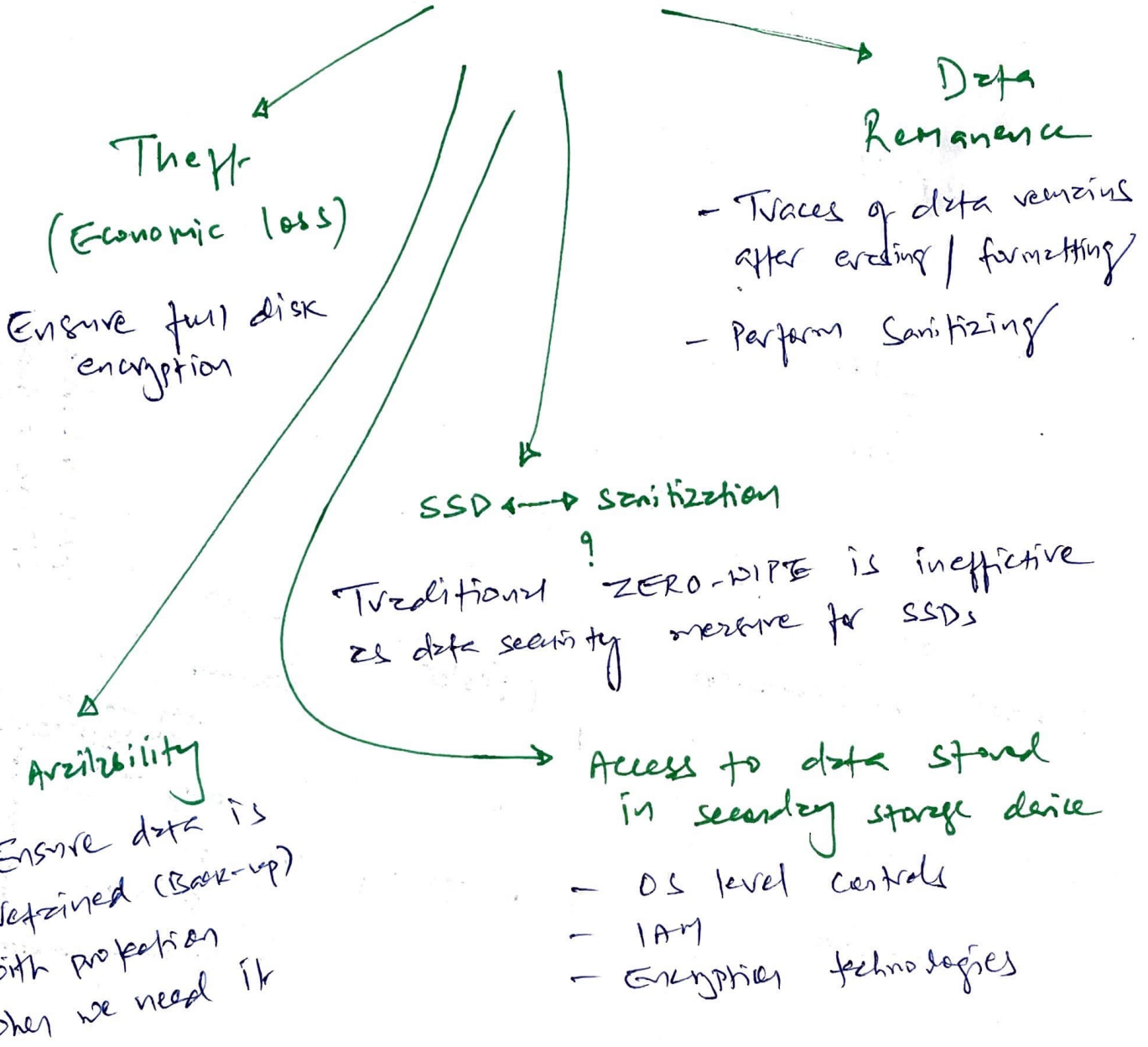
Power-off ⇒ Data lost = Volatile (static / dynamic RAM)

↓
No Data loss

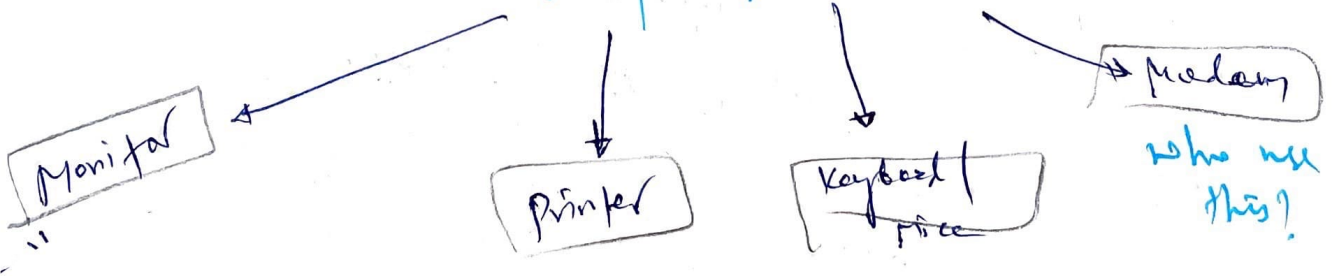
=
Non volatile (magnetic type)

↳ Storage media Security

3 concerns for secondary storage devices



↳ Input / output devices



FIRMWARE

microcode

A software stored in ROM chip.

2 Types

BIOS on motherboard

Basic Input/output system =
instruction to load OS
from disk & start PC

- stored in EEPROM

- Flashing the BIOS =
Process of updating
BIOS

Phishing

Malicious code embedded
into BIOS firmware =
can remotely control the
device

Internal and
External device
firmware

- Mini OS as firmware
in device such as
Printer, Canon C100

- stored in EEPROM =
easy to upgrade

Remember

BIOS =
EEPROM

2011 UEFI replaced BIOS

Unified Extensible Firmware Interface,
advanced interface b/w hardware & OS

CLIENT-BASED SYSTEMS

client-side ATTACK ON

Example

Malicious website



2 things

APPLETS

Mortgage calculator

- Code objects sent from server to client to perform some action
- They are not gone, browsers still support = security risk

- Any data temporarily stored on client system for future use / reuse.

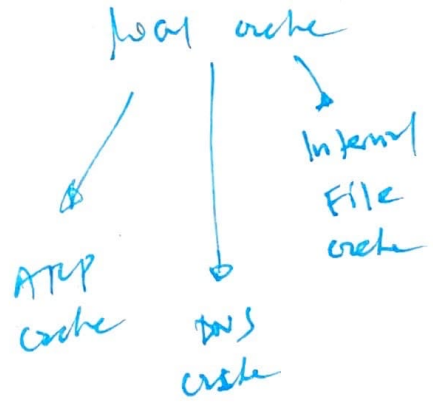
LOCAL CACHE

How?

Remote system send code to local system for execution, if (code) could be trojan.

Mortgage calculator

- Financial data from user could capture & send to server without conscent.



P.T.O X 2
Positioning

Applet types → P.T.O

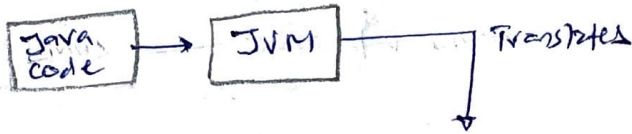
Applet types

JAVA Applet

Problem → Need multiple compilers to produce different versions of single application for each platform it must support.

Soln = Java virtual machine (JVM)

Each system that runs Java code downloads the version of JVM supported by its OS.



To format executable by specific system

JVM

Benefit: code can be shared b/w different OS without modification

Container is one step ahead — deal with only one OS!

Sandbox concept ← To Address security

Java code isolate from OS & strict controls perform on what resources ~~at~~ those objects can access.

Active X Controls

Bed = outdated

- Microsoft's answer to Sun's Java applet

① Active X = Proprietary only runs on Microsoft Browser

② Active X of sandbox → They have full access

Be careful before downloading Active X & executing them.

contd. - Local cache

ARP cache poisoning



DNS Cache Poisoning

Mitm: man-in-the-middle attack

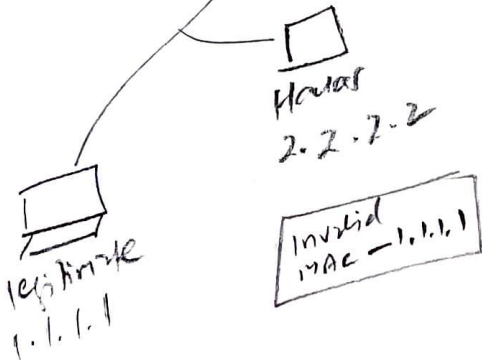
Kali's Ettercap

Poisons dynamic ARP entries (10 min)

Static ARP Poisons = permanent, even after reboots

Kaminsky DNS vulnerability

A → who is 1.1.1.1



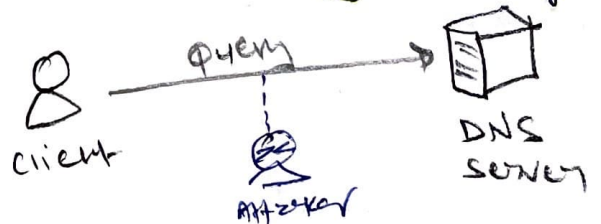
① Static Host file
- Attacker tamper / inject false info b/w FQDN & IP Address relation for permanent damage

② Authorised DNS server attacks
- Primary record of FQDN is stored on primary authoritative DNS server = propagated to Internet
Not effective - Try this

③ Caching DNS servers
- Most ISP cache contents

④ Alternate DNS IP
- clients get wrong DNS IP addresses

⑤ DNS query spoofing



- similar to MITM attack

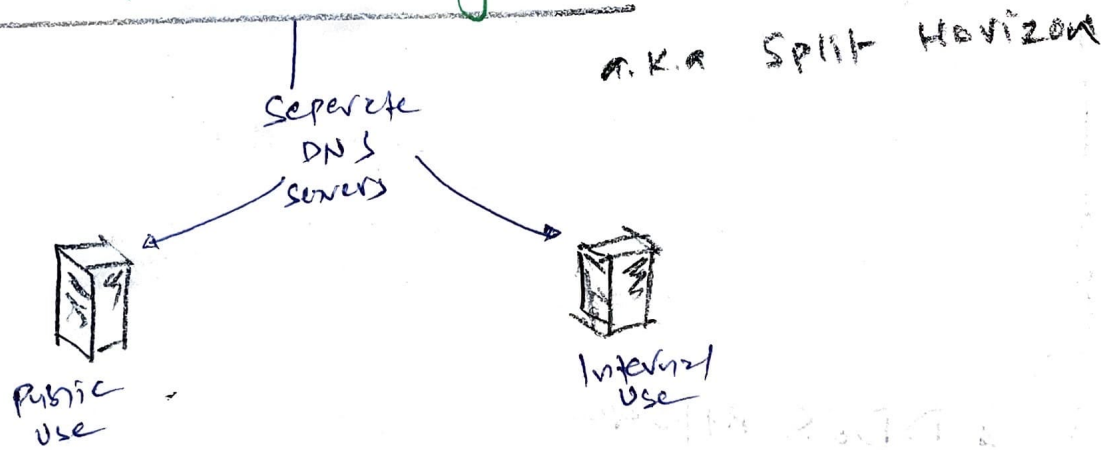
Temporary Internet Files (Internet file cache)

- Website content / downloads
- **Split Response Attack**: client download files from unintended web page.
- **Mobile code scripting Attack**: plant false content in cache

Keep OS +
apps
with
latest
patches

solution

Use SPLIT-DNS system (Split-Brain)



SERVER BASED SYSTEMS

Focus / concern is

DATA FLOW CONTROL

Movement of data b/w process, devices, across the network / channels.

Transmission of Data

- latency
- throughput

considers

Info. protection

- CIA

System is not loaded with too much traffic

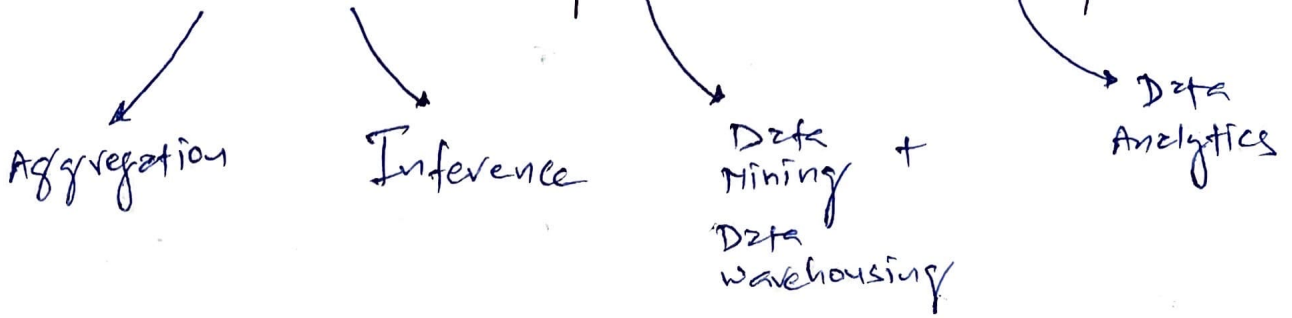
Load Balancer

DDoS Attack

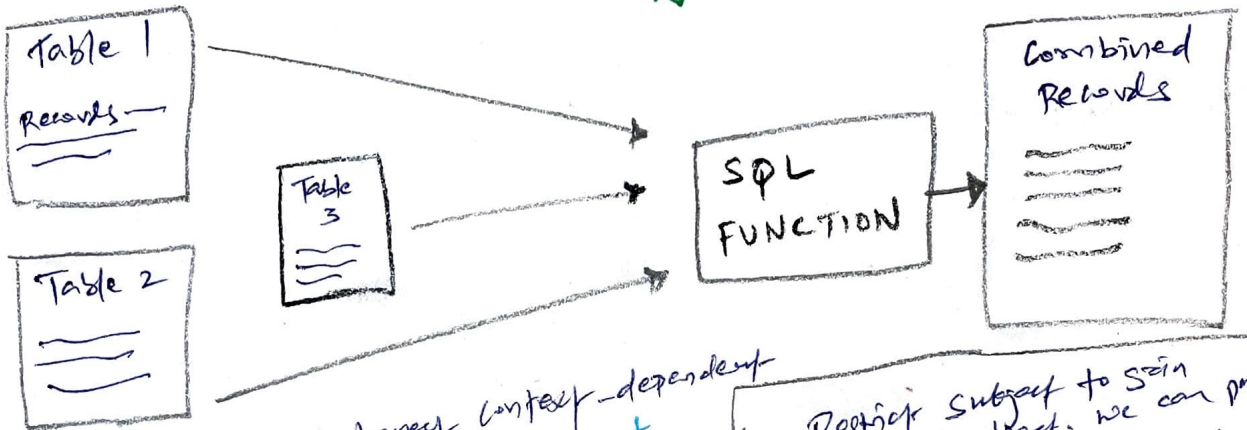
Severe detriment of data flow control

--- Protection mechanisms (ch: 12 to 17)

DATABASE SYSTEM SECURITY.



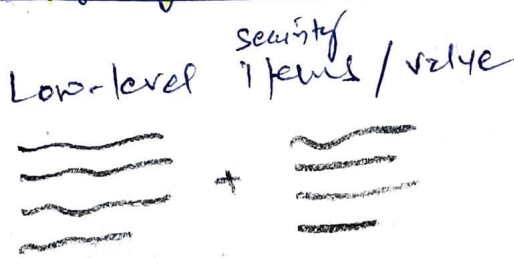
1 AGGREGATION → Combines records from one or more tables to produce useful information.



Incomplete context-dependent Access control

Aggregation Attacks

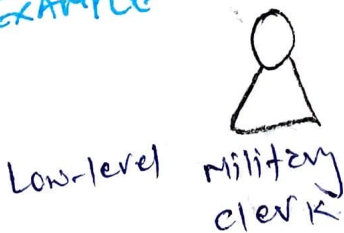
Prevention: Restrict subject to gain access to object. we can put object into higher clearance level so lower clearance subject can't access



Poly/morphism

Creates higher security level/value

EXAMPLE



Equipment Inventory of transfer generally from base to base

Aggregated function =

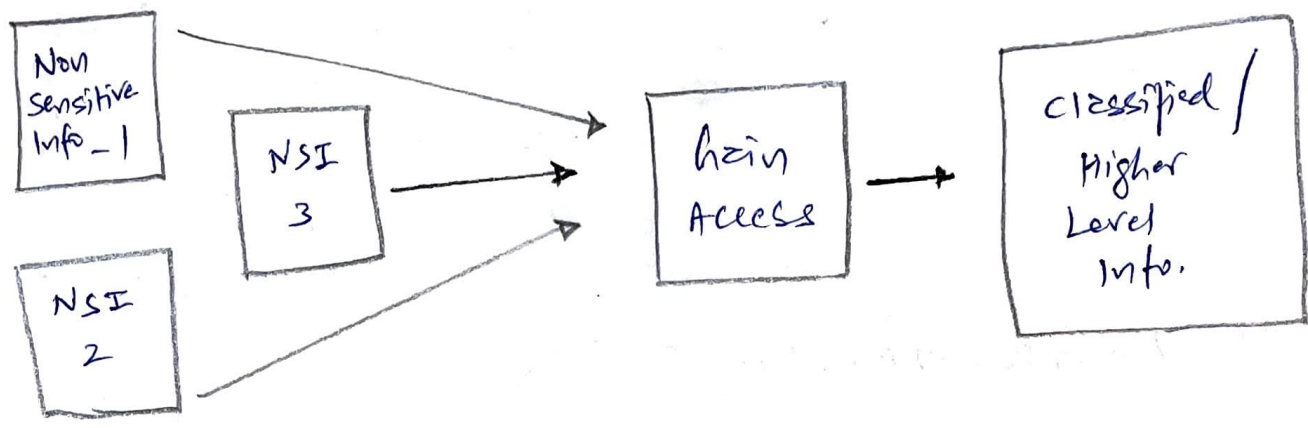
clerk knows how many troops are assigned to each bus station.

SENSITIVE INFO FOR ENEMY

Requires restricted database access.

2 INFERENCE

is intended result of Aggregation



EXAMPLE



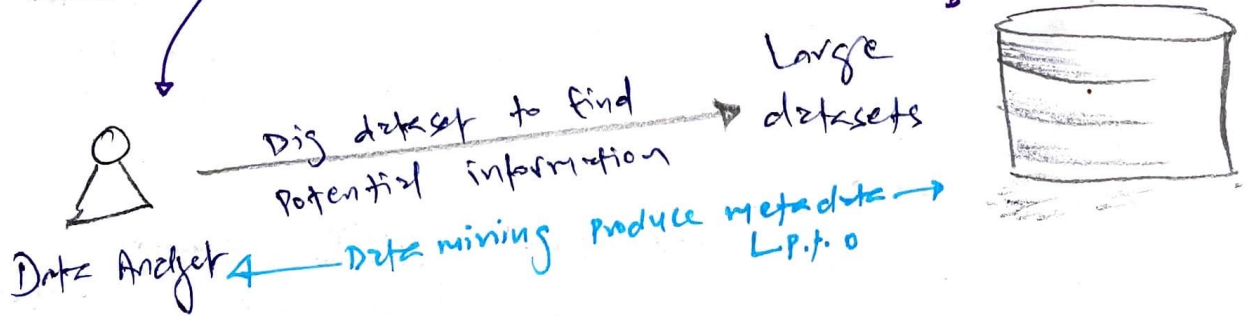
Polyinstantiation
Defense / Protection

- cell suppression
- Database partitioning, noise & perturbation
- Intentional blurring of data
- least privilege

P.t.o
End
bigger picture

vulnerable to attack - 1
& - 2

3 DATA MINING & DATA WAREHOUSING



- critical information about data, usage, type, relationship, format & sources.

- Metadata

- Data about data. Info. about data.
- Can be subset / superset of large dataset.

Example - Security Incident Report

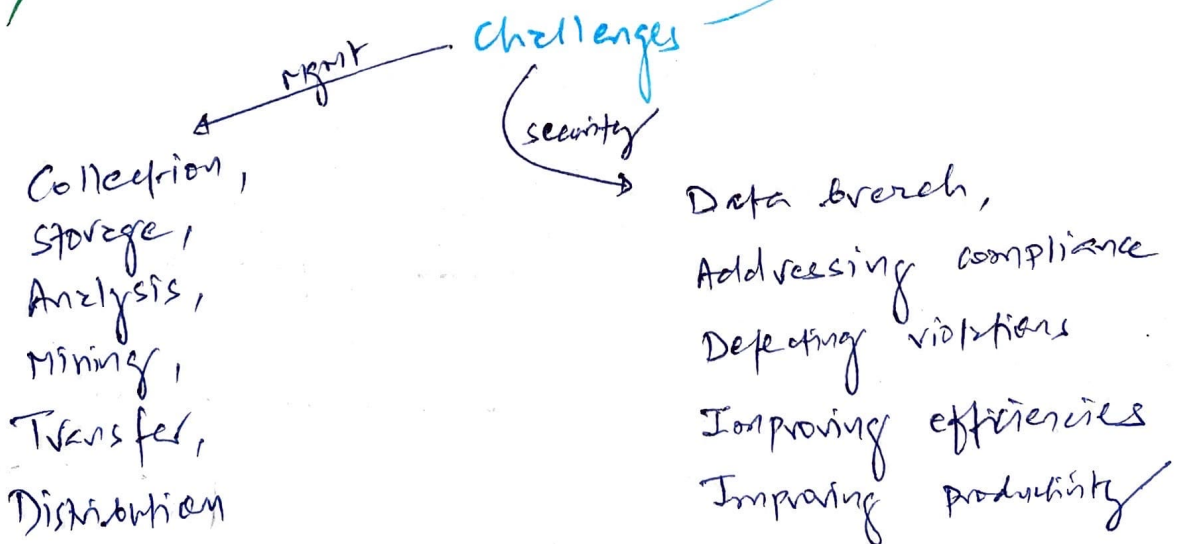
Metadata extracted from large dataset of audit logs

- Metadata is sensitive = security concern

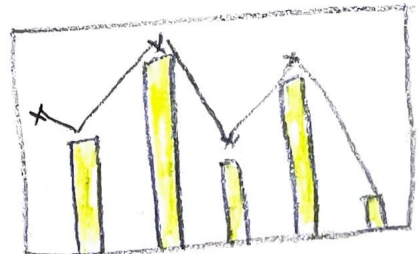
↓
stored in secure container
called

DATA MART

4 DATA ANALYTICS ----- BIG DATA



Extract useful information from bulky dataset that make sense / Insights for business

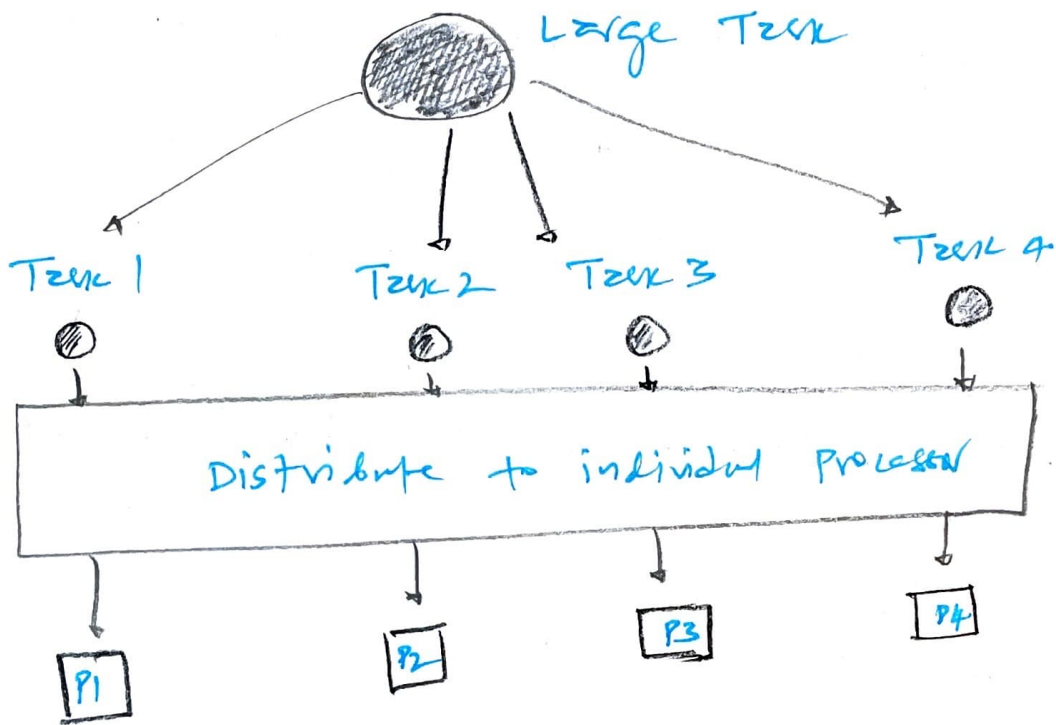


LARGE-SCALE PARALLEL DATA SYSTEMS

Parallel Computing

— multiprocessing

Computational system designed to perform numerous calculations simultaneously.



Multiprocessing Divisions

Symmetric multiprocessing (SMP)

- Processors share same OS & memory
- collection of processors work on same task, code or project.

Asymmetric Multiprocessing (AMP)

- Processors & OS are independent of each other
- Each processor has its own OS & task

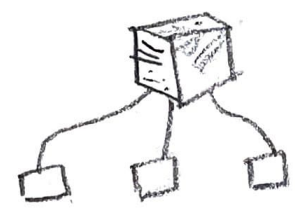
— Under AMP, processors are configured to execute specific task/code, it's called Affinity in some circumstances.

DISTRIBUTED SYSTEM AND ENDPOINT SECURITY

Evolution of Computing



Host / Terminal Model



Client-Server Model

Distributed Architecture

- Prove to vulnerabilities in monolithic host / terminal system

- Communication equipment can be unwanted points of entry to distributed environment.
 Router / switch / Modem

- User = threat if download malicious code / Trojan horse

- Data on machine = risk if not properly backed-up.

Security concern

Process + storage distributed on multiple clients & servers = Everything must be secured.

Virus spreads faster in distributed compare to monolithic architecture

To Safeguard Distributed Environment means

